

TCP over ATM using ABR, UBR, and GFR Services and QoS over IP Issues

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- ❑ Why ATM?
- ❑ ABR Vs UBR
- ❑ TCP/IP over UBR
- ❑ TCP/IP over GFR
- ❑ QoS over IP: IntServ, DiffServ, MPLS

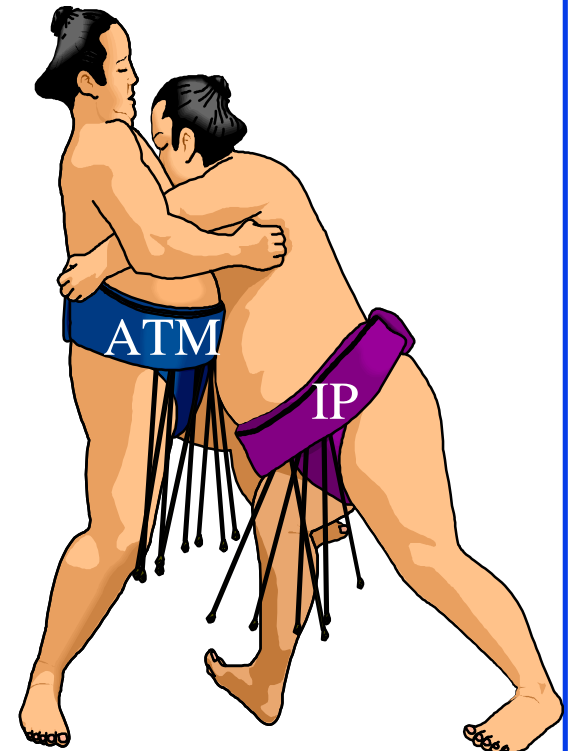
Ref: For detailed studies, see

<http://www.cis.ohio-state.edu/~jain/>

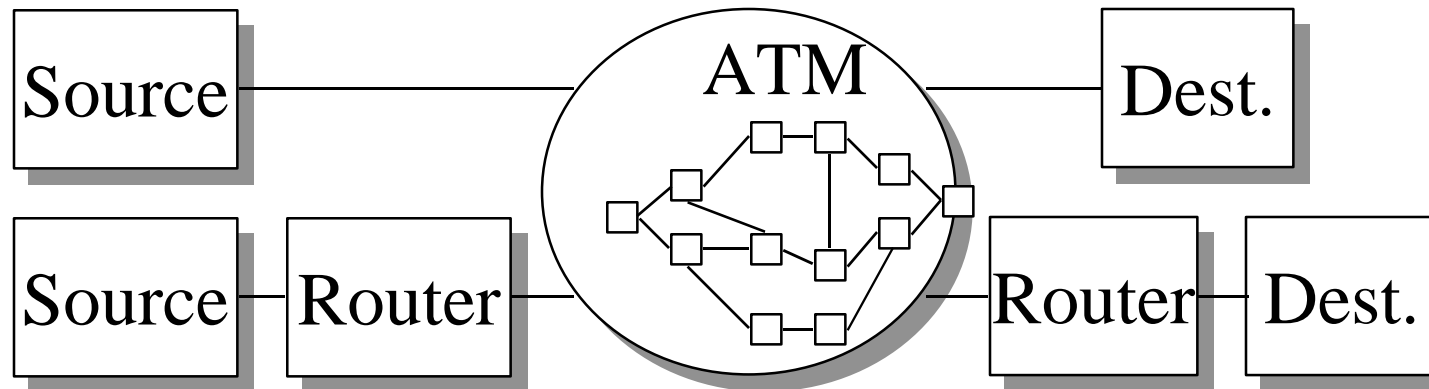
Why ATM?

ATM vs IP: Key Distinctions

1. Traffic Management:
Explicit Rate vs Loss based
2. Signaling: Coming to IP in the form of RSVP
3. QoS: PNNI routing, Service categories. Integrated/Differentiated services
4. Switching: Coming to IP as MPLS
5. Cells: Fixed size or small size is not important



ABR vs UBR



ABR

Queue in the source
Pushes congestion to edges
If ATM not end-to-end:
intelligent Q mgmt in
routers
Works for all protocols

UBR

Queue in the network
No backpressure
Same end-to-end or backbone
Works with TCP

Improving TCP over UBR

End-System Policies

		No FRR	FRR	New Reno	SACK + New Reno	
		Switch Policies	No EPD			
Plain EPD						
EPD	Selective Drop					
	Fair Buffer Allocation					

Policies: Results

- ❑ In LANs, switch improvements (PPD, EPD, SD, FBA) have more impact than end-system improvements (Slow start, FRR, New Reno, SACK). Different variations of increase/decrease have little impact due to small window sizes.
- ❑ In large bandwidth-delay networks, end-system improvements have more impact than switch-based improvements
- ❑ FRR hurts in large bandwidth-delay networks.

Policies (Continued)

- ❑ Fairness depends upon the switch drop policies and not on end-system policies
- ❑ In large bandwidth-delay networks:
 - SACK helps significantly
 - Fairness is not affected by SACK
- ❑ In LANs:
 - Previously retransmitted holes may have to be retransmitted on a timeout
⇒ SACK can hurt under extreme congestion.

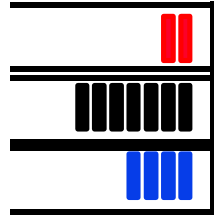
Guaranteed Rate Service

- Guaranteed Rate (GR): Reserve a small fraction of bandwidth for UBR class.

GR	GFR
per-class reservation	per-VC reservation
per-class scheduling	per-VC accounting/scheduling
No new signaling	Need new signaling
Can be done now	In TM4+

- For WANs, the effect of reserving 10% bandwidth for UBR is more than that obtained by EPD, SD, or FBA. For LANs, guaranteed rate is not so helpful. Drop policies are more important.

GFR: Results



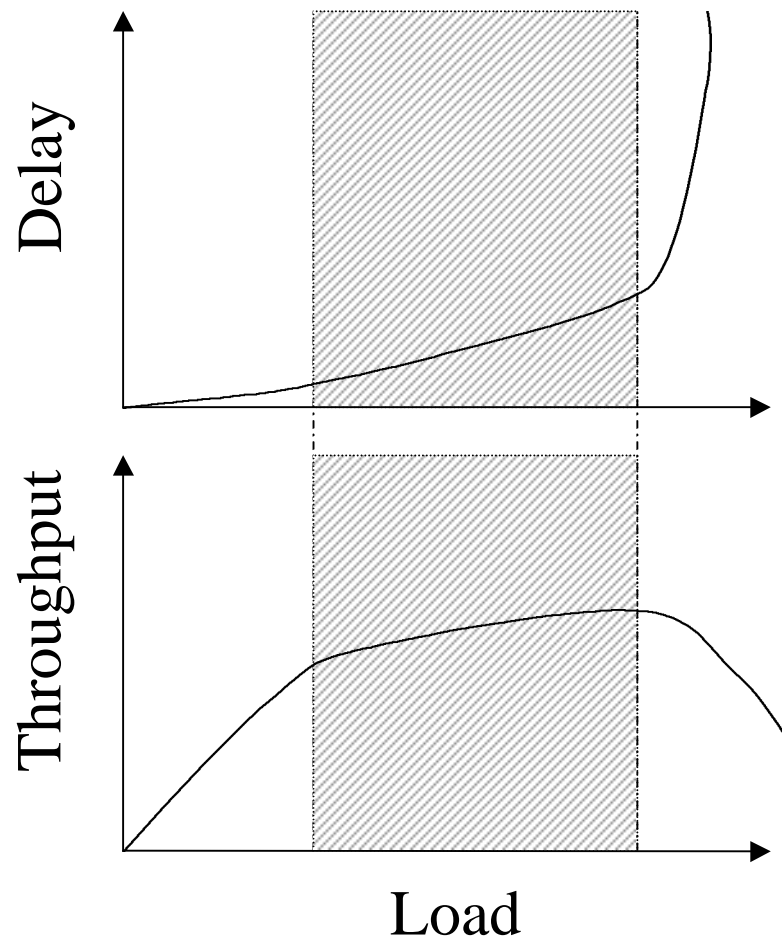
Per-VC Q



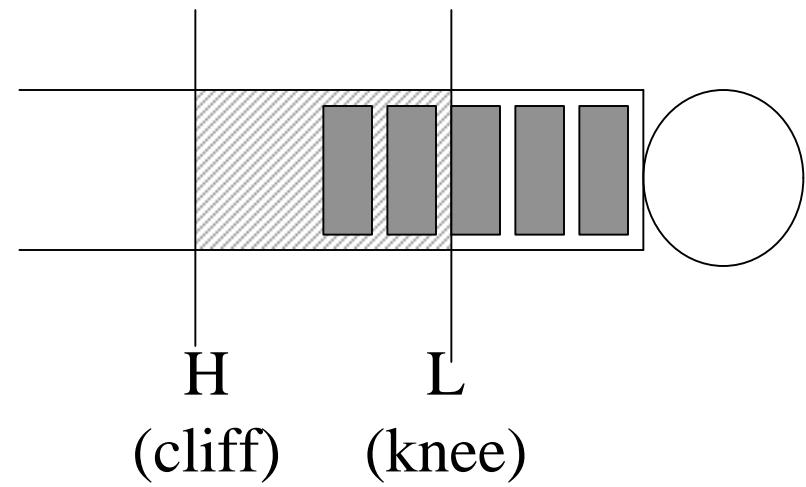
Single FIFO

- ❑ Per-VC queuing and scheduling is sufficient for per-VC MCR.
- ❑ FBA and proper scheduling is sufficient for fair allocation of excess bandwidth
- ❑ Questions:
 - How and when can we provide MCR guarantee with FIFO?
 - What if each VC contains multiple TCP flows?

Differential Fair Buffer Allocation

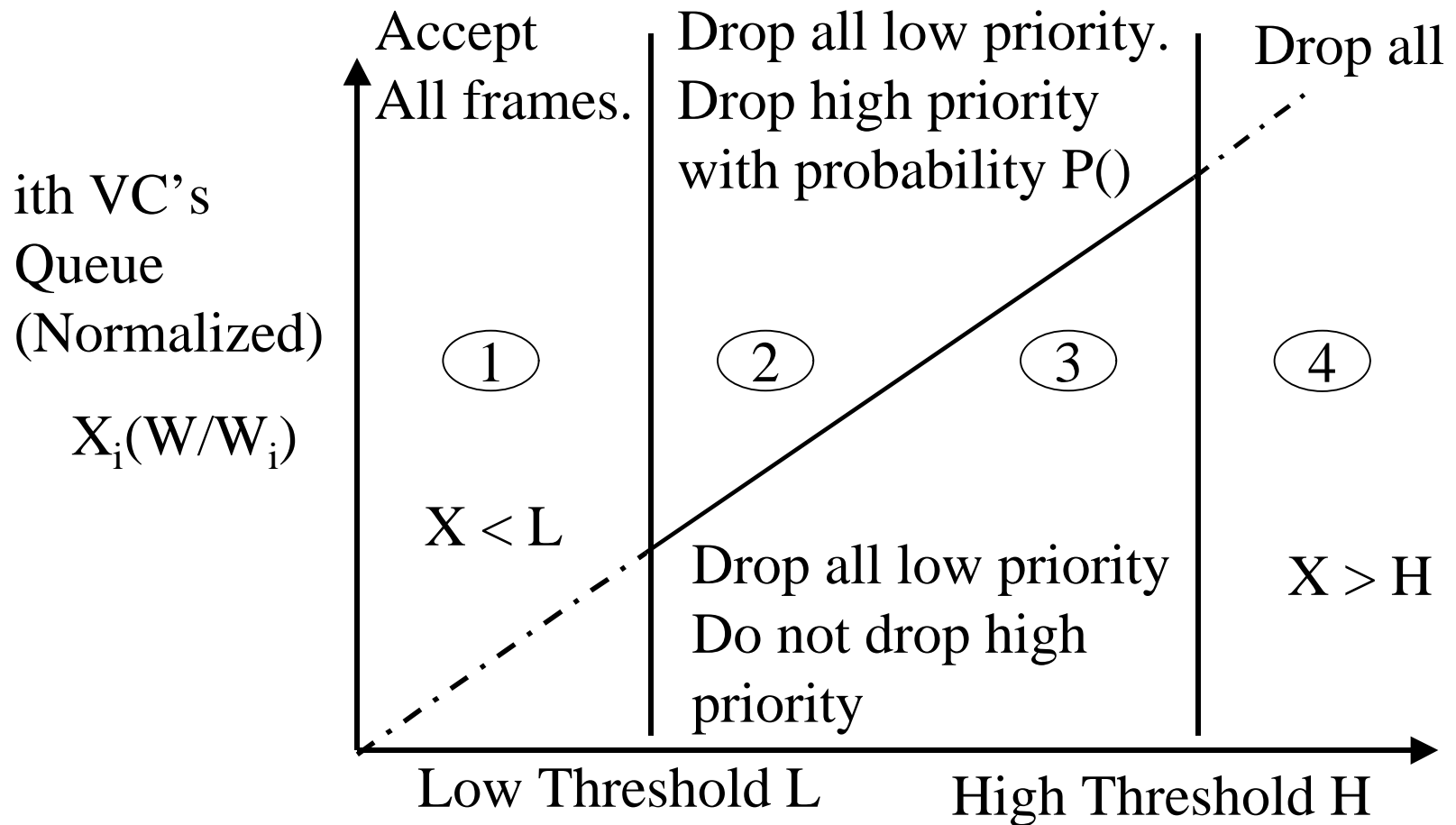


Buffer occupancy (X)



 Desired operating region

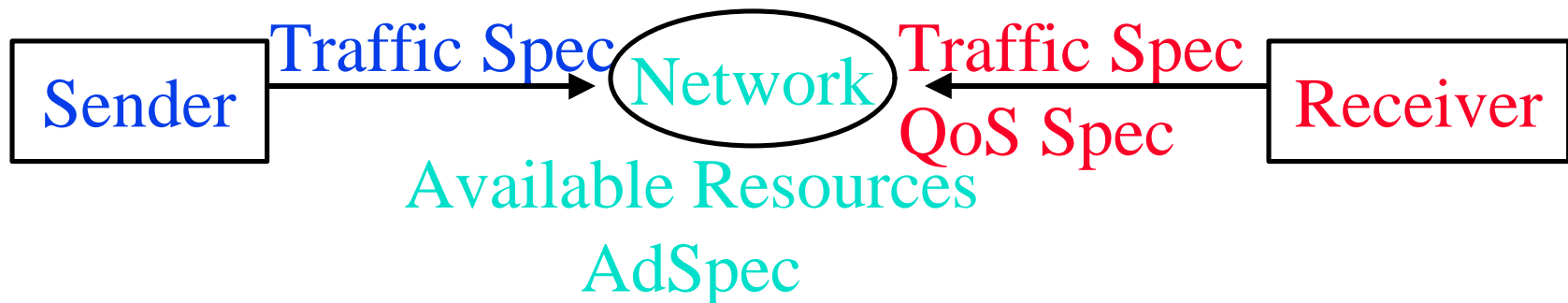
DFBA (contd.)



TCP Rate $D \propto \frac{MSS}{RTT \times \sqrt{P(drop)}}$

Integrated Services and RSVP

- ❑ Best Effort Service: Like UBR.
- ❑ Controlled-Load Service: Performance as good as in an unloaded datagram network. No quantitative assurances. Like nrt-VBR or UBR w MCR
- ❑ Guaranteed Service: Like CBR or rt-VBR
 - Firm bound on data throughput and delay.
 - Is not always implementable, e.g., Shared Ethernet.
- ❑ Resource ReSerVation Protocol: Signaling protocol



Problems with RSVP and Integrated Services

- ❑ Complexity: Packet classification, Scheduling
- ❑ Scalable in number of receivers per flow but
Per-Flow State: $O(n)$ \Rightarrow Not scalable with # of flows.
Number of flows in the backbone may be large.
 \Rightarrow Suitable for small private networks
- ❑ Need a concept of “Virtual Paths” or aggregated flow groups for the backbone
- ❑ Need policy controls: Who can make reservations?
Support for accounting and security.
- ❑ RSVP does not have negotiation and backtracking

Differentiated Services

Ver	Hdr Len	Precedence	ToS	Unused	Tot Len
4b	4b	3b	4b	1b	16b

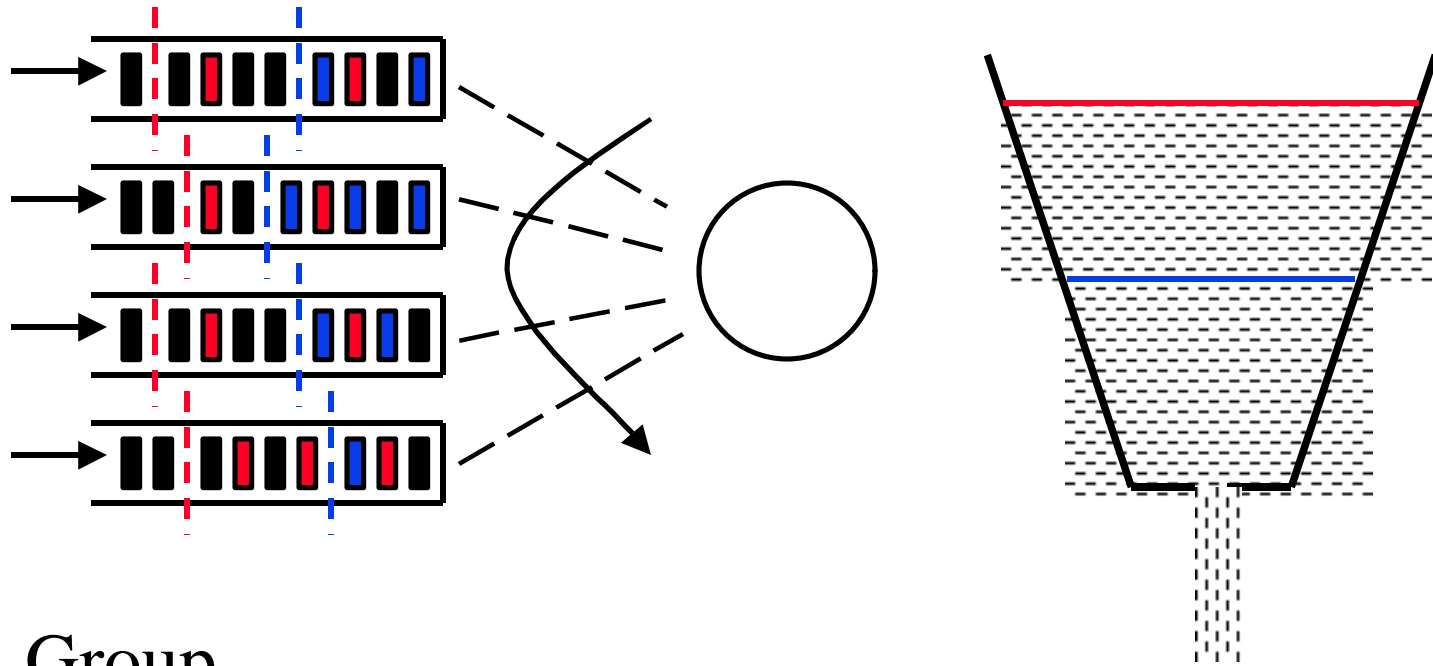
- ❑ IPv4: 3-bit precedence + 4-bit ToS
- ❑ Many vendors use IP precedence bits but the service varies \Rightarrow Need a standard \Rightarrow Differentiated Services
- ❑ DS working group formed February 1998
- ❑ Charter: Define ds byte (IPv4 ToS field)
- ❑ Per-Hop Behavior: Externally Observable Forwarding Behavior, e.g., x% of link bandwidth, or priority



Expedited Forwarding

- ❑ Also known as “Premium Service”
- ❑ Virtual leased line
- ❑ Similar to CBR
- ❑ Guaranteed minimum service rate
- ❑ Policed: Arrival rate $<$ Minimum Service Rate
- ❑ Not affected by other data PHBs
 - ⇒ Highest data priority (if priority queueing)

Assured Forwarding



- ❑ PHB Group
- ❑ Four Classes: Decreasing weights in WFR/WFQ
- ❑ Three drop preference per class
(one rate and two bucket sizes)

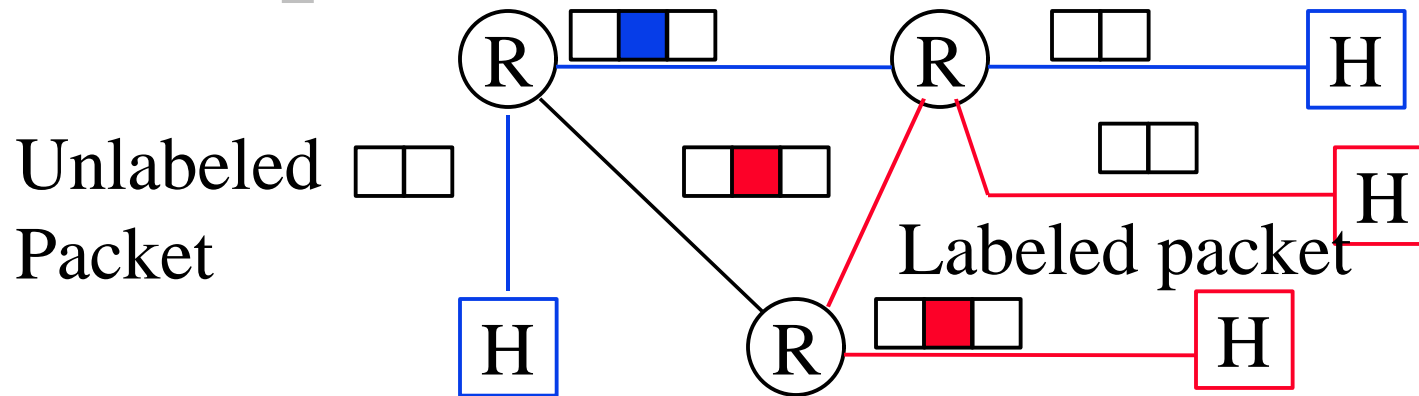
Problems with DiffServ

- ❑ per-hop \Rightarrow Need at every hop
One non-DiffServ hop can spoil all QoS
- ❑ End-to-end $\neq \Sigma$ per-Hop
Designing end-to-end services with weighted guarantees at individual hops is difficult.
Only EF will work.
- ❑ QoS is for the aggregate not micro-flows.
Not intended/useful for end users. Only ISPs.
 - Large number of short flows are better handled by aggregates.

DiffServ Problems (Cont)

- Long flows (voice and video sessions) need per-flow guarantees.
- High-bandwidth flows (1 Mbps video) need per-flow guarantees.
- All IETF approaches are open loop control \Rightarrow Drop.
Closed loop control \Rightarrow Wait at source
Data prefers waiting \Rightarrow Feedback
- Guarantees \Rightarrow Stability of paths
 \Rightarrow Connections (hard or soft)
Need route pinning or connections.

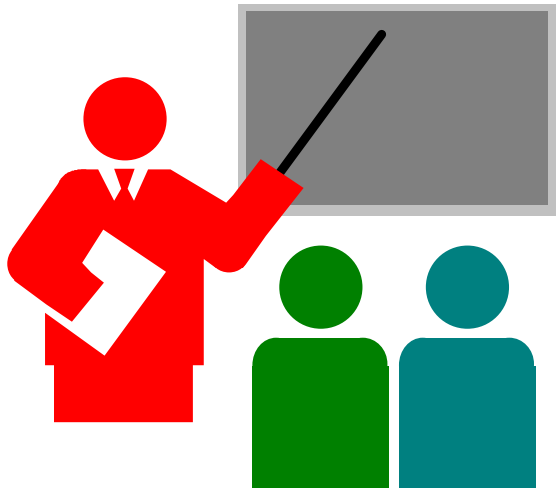
Multiprotocol Label Switching



- ❑ Entry “label switch router (LSR)” attaches a label to the packet based on the route
- ❑ Other LSRs switch packets based on labels. Do not need to look inside \Rightarrow Fast.
- ❑ Labels have local significance \Rightarrow Different label at each hop (similar to VC #)
- ❑ Exit LSR strips off the label

Traffic Engineering Using MPLS

- ❑ Traffic Engineering = Performance Optimization
= Efficient resource allocation, Path splitting
⇒ Maximum throughput, Min delay, min loss
⇒ Quality of service
- ❑ In MPLS networks: “Traffic Trunks” = SVCs
Traffic trunks are routable entities like VCs
- ❑ Multiple trunks can be used in parallel to the same egress.
- ❑ Each traffic trunk can have a set of associated characteristics, e.g., priority, preemption, policing, overbooking



Summary

- ❑ Traffic management distinguishes ATM from its competition
- ❑ ABR pushes congestion to edges.
UBR+ may be OK for LANs but not for large bandwidth-delay paths.
- ❑ Reserving a small fraction of bandwidth for the entire UBR class improves its performance considerably.
- ❑ It may be possible to do GFR with FIFO

Summary

- ❑ Multiple drop preferences does not help data (TCP) or Voice/Video
- ❑ Voice/video need multiple leaky bucket rates for layered/scalable coding.
- ❑ Need additivity or mathematical aggregatability. CBR (EF) should be the first step for IP.
- ❑ Excess allocation is useful with closed loop. Network/application dynamics \Rightarrow Need closed loop