

99-0045

Throughput Fairness Index: An Explanation

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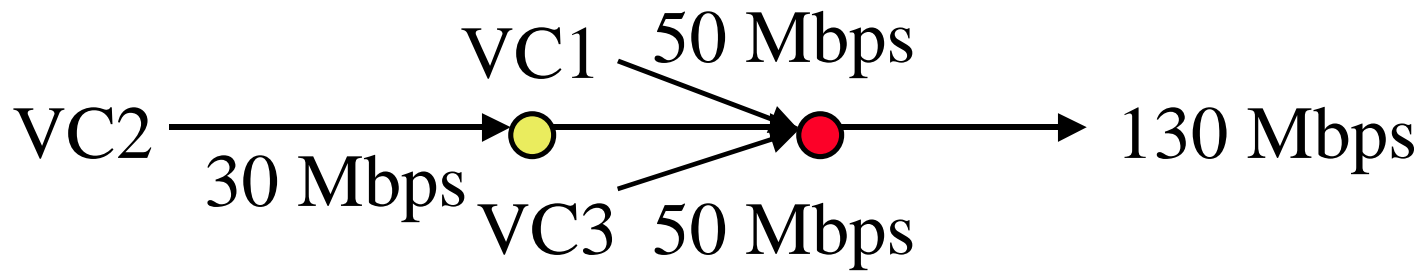


- ❑ Index of fairness
- ❑ Why is it better than others?

References:

[1] R. Jain, W. Hawe, D. Chiu, “A Quantitative measure of fairness and discrimination for resource allocation in Shared Computer Systems,” DEC-TR-301, September 26, 1984, <http://www.cis.ohio-state.edu/~jain/papers/fairness.htm>

Fairness



- ❑ Simple Definition: Equal share of bottleneck
Problem: Some VC's may be bottlenecked elsewhere
- ❑ Solution: Use any of the TM4.0 fairness criteria to define optimal allocations and use a fairness index to quantify the fairness.
- ❑ Question: A scheme gives 50, 30, 50 Mbps when the optimal is 50, 10, 10 Mbps
How fair is it? 67% ? 90% ?

Proposal

- ❑ Measured Throughput: (T_1, T_2, \dots, T_n)
- ❑ Use any criterion (e.g., max-min optimality) to find the Fair Throughput (O_1, O_2, \dots, O_n)
- ❑ Normalized Throughput: $x_i = T_i/O_i$

$$\text{Fairness Index} = \frac{(\sum x_i)^2}{n\sum x_i^2}$$

Example: 50/50, 30/10, 50/10 \Rightarrow 1, 3, 5

$$\text{Fairness Index} = \frac{(1+3+5)^2}{3(1^2+3^2+5^2)} = \frac{9^2}{3(1+9+25)} = 0.81$$

Other Fairness Indices

q Variance:

$$\text{Mean } \mu = (1+3+5)/3 = 3$$

$$\text{Variance } \sigma^2 = 1/(n-1)\Sigma(x_i-\mu)^2 = 4$$

q Coefficient of Variation: Standard deviation $\sigma = 2$

$$\text{COV} = \sigma/\mu = 0.667$$

q Min-Max Ratio: $\text{Min}\{x_i\}/\text{Max}\{x_i\} = 1/5 = 0.2$

q Find the normalized distance from the optimal

$$\begin{aligned} \text{Norm. Dist.} &= \frac{[\Sigma (T_i - O_i)^2]^{1/2}}{[\Sigma O_i^2]^{1/2}} \\ &= \frac{[(50-50)^2 + (30-10)^2 + (50-10)^2]^{1/2}}{[50^2 + 10^2 + 10^2]^{1/2}} = 0.86 \end{aligned}$$

Fairness Index: Properties

- ❑ Applicable for **any number of VCs**, even $n=2$
Strictly speaking, variance not defined for small n .
- ❑ **Scale independent.**
Variance (Throughput) = $10 \text{ Mbps}^2 = 10^7 \text{ kbps}^2$
Standard deviation (Throughput) = $10 \text{ Mbps} = 10^4 \text{ kbps}$
- ❑ **Bounded** between 0 and 1 or 0 and 100%
Variance, standard deviation, and Relative distance are not bounded.
- ❑ **Direct relationship:** Higher index \Downarrow More Fair
Higher variance \Downarrow Less fair
- ❑ **Continuous.** Min/max is not continuous.

Fairness Index: Properties

Intuitive:

□ For $(1, 0, 1)$ Index = $2/3$

□ For $x_i = 1, i=1,2,3,\dots,k$
= 0 otherwise

Fairness Index = k/n

□ If 80% of the users are treated fairly and 20% are starved, index = 80%

□ If $y\%$ of the users are treated fairly and $(100-y)\%$ are starved, Fairness index = $y\%$

Relationship to Other Indices

□ Fairness Index = $E[x]^2/E[x^2] = 1/(1+COV^2)$

Transformation

- Makes index bounded between 0 and 1,
- Gives a direct positive relationship between the index and fairness
- Makes it intuitive: $y\%$ if fair to $y\%$ only.

User Perception of Fairness

- Fairness Index = $(\sum x_i)^2 / (n \sum x_i^2)$
= $(1/n) \sum x_i / x_f$

Where $x_f = \sum x_i^2 / \sum x_i =$ Fair allocation mark

- i^{th} User perception of fairness = x_i / x_f

Example: 2 Mbps to first 10 users, 0 to other 90 users

$$x_i = 2, i=1,2,\dots,10$$

$$x_i = 0, i=11,12,\dots,100$$

Fair Allocation Mark $x_f = 2$

First 10 users are 100% happy and other 90 users are 0% happy. Average fairness = 10%

Properties of Fairness Index

1. If δx resource is taken away from k th user and given to j th user, the fairness

- Increases iff $\delta x < x_k - x_j$
- Remains the same iff $\delta x = x_k - x_j$
- Decreases iff $\delta x > x_k - x_j$

2a. If each user is given an additional resource δx , the index goes up if it is less than 1.

$$f(x_1 + c, x_2 + c, \dots, x_n + c) \geq f(x_1, x_2, \dots, x_n)$$

2b. If a single user j is given extra δx , the index goes up iff j is a discriminated user (and vice versa):

$$f(x_1, x_2, \dots, x_j + \delta x, \dots, x_n) > f(x_1, x_2, \dots, x_n) \text{ iff } x_j < x_f$$

Properties (Cont)

3. If we vary a single user's allocation, the fairness reaches maximum when $x_j = x_g$ where x_g is the fair allocation mark for the remaining $n-1$ users

$$x_g = \frac{\sum_{i \neq j} x_i^2}{\sum_{i \neq j} x_i}$$

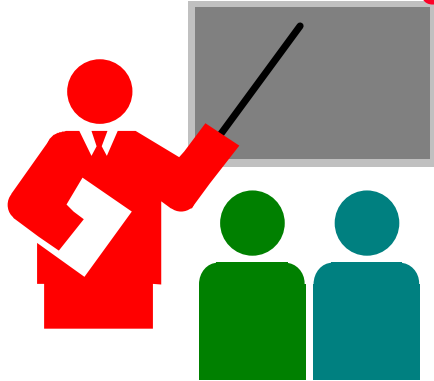
Other Similar Functions

- $F(x) = (\sum x_i)^k / (n \sum x_i^k)$
- If you starve 90% of the users and give all resources to 10% of the users:

$$(\sum x_i)^k / (n \sum x_i^k) = 0.1^{k-1}$$

The index is 0.1 or 10% only for $k=2$.

Summary



$$\text{Fairness Index} = \frac{(\sum x_i)^2}{n \sum x_i^2}$$

- ❑ Independent of scale
- ❑ Continuous
- ❑ Applies to any number of users
- ❑ Bounded between 0 and 1
- ❑ Has a intuitive relationship with user perception